

# 1 Bit Full Adder Cell for Reducing Low Leakage Current in Nano Meter Technology

Madhuri.Sada<sup>1</sup>, A.Srinivasulu<sup>2</sup>, C.Md.Aslam<sup>3</sup>,

<sup>1</sup>M. Tech Student, <sup>2</sup>M.Tech, Assistant Professor, Department of Electronics (E.C.E), Vaagdevi Institute of Tech & Science, Peddasettipalli(Village), Proddatur(M), Kadapa(Dist), Andhra Pradesh A.P, India.

<sup>3</sup>Associate Professor, Head of the Department of Electronics (E.C.E), Vaagdevi Institute of Tech & Science, Peddasettipalli(Village), Proddatur(M), Kadapa(Dist), Andhra Pradesh A.P, India.

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**Abstract**— as technology scales into the nanometer regime leakage current, active power, delay and area are becoming important metric for the analysis and design of complex arithmetic logic circuits. In this paper, low leakage 1bit full adder cells are proposed for mobile applications and a novel technique and gated-diffusion input GDI technique have been introduced for further reduction in power. We introduced a new transistor resizing approach for 1bit full adder cells to determine the optimal sleep transistor size which reduce the leakage power and area. The simulation results depicts that the propose design also leads to efficient 1-bit full adder cells in terms of standby leakage power, active power. We have performed simulations using Cadence Spectra 90nm standard CMOS technology at room temperature with supply voltage of 1V.

**Keywords**— Low leakage power; Active power; Sleep transistor and Adder cell.

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## I. INTRODUCTION

Adders are heart of computational circuits and many complex arithmetic circuits are based on the addition. The vast use of this operation in Arithmetic functions attracts a lot of researcher's attention to adder for mobile applications. In recent years, several variants of different logic styles have been proposed to implement 1-bit adder cells. These adder cells commonly aimed to reduce power consumption and increase speed. These studies have also investigated different approaches realizing adders using CMOS technology. For mobile applications, designers have to work within a very tight leakage power specification in order to meet product battery life and package cost objectives. The designer's concern for the level of leakage current is not related to ensuring correct circuit operation, but is related to minimize power dissipation. For portable electronic devices this equates to maximizing battery life. For example, mobile phones need to be powered for extended periods (known as standby mode, during which the phone is able to receive an incoming call), but are fully active for much shorter periods (known as talk or active mode, while making a call). When an electronic device such as a mobile phone is in standby mode, certain portions of the circuitry within the electronic device, which are active when the phone is in talk mode, are shut down. These circuits, however, still have leakage currents running through them, even though they have been de-activated. Even if the leakage current is much smaller than the normal operating current of the circuit. The leakage current depletes the battery charge over the relatively long standby time, whereas the operating current during talk time only depletes the battery charge over the relatively short talk time. As a result, the leakage current has a disproportional effect on total battery life. This is why building low leakage adder cells for mobile applications are of great interest.

To summarize, some performance criteria are considered in the design and evaluation of adder cells, such as leakage power, active power, ground bounce noise, area, noise margin and robustness with respect to voltage and transistor scaling as well as varying process and compatibility with surrounding circuitries.

Shortening the gate length of a transistor increases its power consumption due to the increased leakage current between the transistors source and drain when no signal voltage is applied at the gate. In addition to the sub threshold leakage current, gate tunneling current also increases due to the scaling of gate oxide thickness. Each new technology generations results nearly a 30x increase in gate leakage. The leakage power is expected to reach more than 50% of total power in sub 100nm technology generation. Hence, it has become extremely important to develop design techniques to reduce static power dissipation during periods of inactivity. The power reduction must be achieved without trading-off performance which makes it harder to reduce leakage during normal (runtime) operation. On the other hand, there are several techniques to reduce leakage power.

Power gating is one such well known technique where a sleep transistor is added between actual ground rail and circuit ground (called virtual ground). This device is turned off in the sleep mode to cut-off the leakage path. It has been shown that this technique provides a substantial reduction in leakage at a minimal impact on performance and further peak of ground bounce noise is possible with proposed novel technique with improved staggered phase damping technique.

This paper focuses on reducing sub threshold leakage power consumption and ground bounce noise. The remainder of this paper is organized as follows. In section II, proposed nano-CMOS full adder circuits, and its equivalent

circuits are discussed. In section III, the performance analysis and simulation results of conventional CMOS full adder cell and proposed circuits are explained.

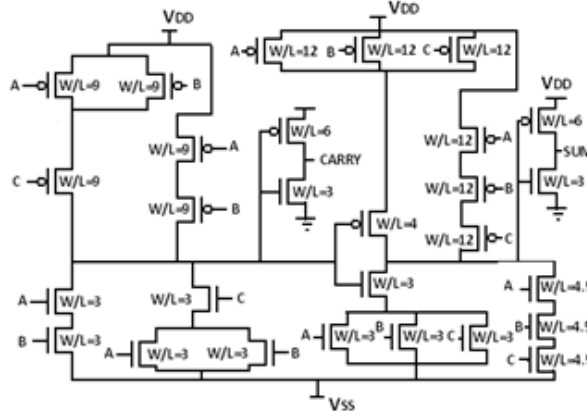
## II. PROPOSED FULL ADDER CIRCUITS

Recently, power dissipation has become an important concern and considerable emphasis is placed on understanding the sources of power and approaches to dealing with power dissipation.

Static logic style gives robustness against noise effects, so automatically provides a reliable operation. Pseudo NMOS and Pass-transistor logic can reduce the number of transistors required to implement a given logic function. But those suffer from static power dissipation. Implementing Multiplexers and XOR based circuits are advantageous when we implement by the pass transistor logic. On the other hand, dynamic logic implementation of complex function requires a small silicon area but charge leakage and charge refreshing are required which reduces the frequency of operation. In general, none of the mentioned styles can compete with CMOS style in robustness and stability.

*Fig. 1* shows the conventional CMOS 28 transistor adder. This is considered as a Base case throughout this paper. All comparisons are done with Base case. The CMOS structure combines PMOS pull up and NMOS pull down networks to produce considered outputs. Transistor sizes are specified as a ratio of Width/Length (W/L). The sizing of transistors plays a key role in static CMOS style. It is observed in the conventional adder circuit that the transistor ratio of PMOS to NMOS is 2 for an inverter and remaining blocks also followed the same ratios when we considered the remaining blocks as an equivalent inverters. This ratio does not give best results with respect to noise margin and standby leakage power when it is simulated in 90nm process. Modified adder circuits with sizing are proposed in Design1 and Design2 targeting the noise margin, and ground bounce noise.

Further, power gating technique is used to reduce the leakage power, where a sleep transistor is connected between actual ground rail and circuit ground. Ground bounce noise is being estimated when the circuits are connected with a sleep transistor. Further, the peak of ground bounce noise is achieved with a proposed novel technique.



*Fig. 1* Conventional CMOS full adder.

Modified sizing's are shown in *Fig. 2* and *Fig. 5* respectively. The smallest transistor considered for 90nm technology has a width of 120nm and a length of 100nm and gives W/L ratio of 1.2. The W/L ratio of NMOS is fixed at 1.2 and W/L of PMOS is 3.8 which is 3.17 times that of NMOS in Design1. The sizing of each block is based on the following assumption.

Base case is considered as individual block as shown in *Fig. 3*. Each block has been treated as an equivalent inverter. The same inverter ratio is maintained on each block. These sizing will reduce the standby leakage current greatly because sub-threshold current is directly proportional to the Width/Length ratio of transistor. On the other hand, these reduced sizes will reduce the area occupied by the circuit. This will reduce the silicon chip area and obviously there will be a reduction in the cost. Modified adder circuit i.e. Design2 shown in *Fig. 5*, the W/L ratio of PMOS is 1.5 times that of W/L ratio of NMOS and each block has been treated as an equivalent inverter. The same inverter size has been maintained on each block as shown in the *Fig. 4*. The goal of this design is to reduce the standby leakage power. Further compared to the Base case and Design1 and ground bounce noise produced when a circuit is connected to sleep transistor. However, there will be a slight variation on the noise margin levels and is almost equal to the Base case.

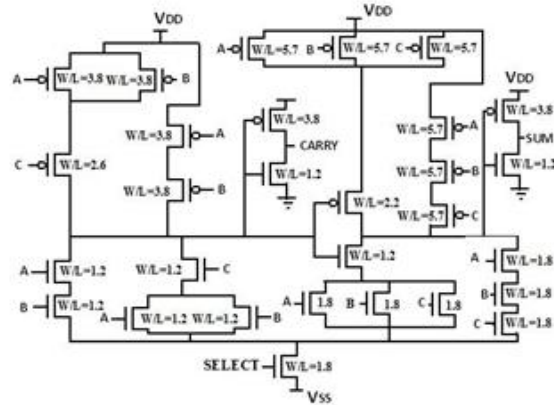


Fig. 2 Proposed full adder (Design1) circuit with sleep transistor

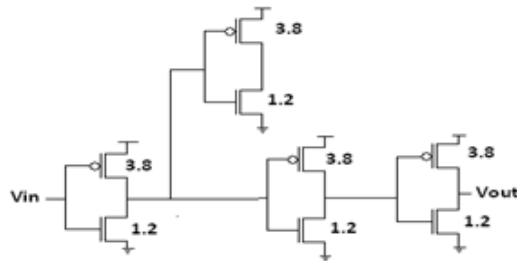


Fig. 3 Equivalent circuit for Design1

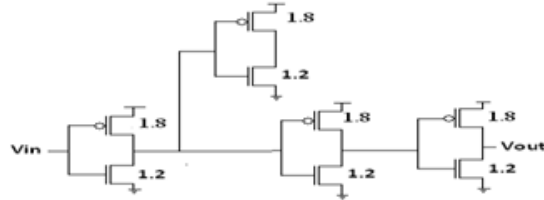


Fig. 4 Equivalent circuit for Design 2.

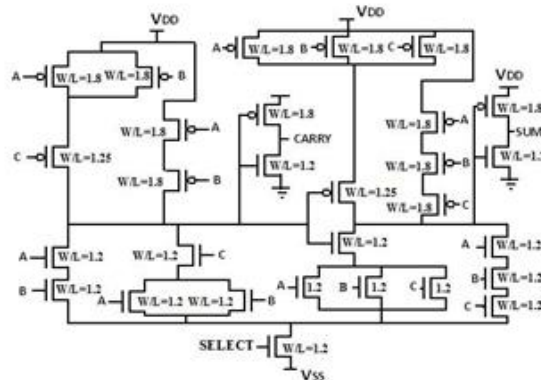


Fig. 5 Proposed 1 bit full adder (Design2) circuit with sleep transistor

### III. PROPOSED NOVEL TECHNIQUE

Fig. 6 shows the proposed novel technique. One bit full adders (Base Case, Design1, design2) have been taken to apply the proposed technique. One-bit full adder considered as two cascaded blocks i.e. carry generation block and sum generation block. Separate sleep transistors are added at the bottom of the blocks. The proposed technique works as follows. For carry part, during stage1 transmission gate is off by giving proper enable signals and at the same time control transistor is turned on to make the sleep transistor working as a

diode. The stored charge in carry generation block is discharged through sleep transistor. The drain current of the ST1 during this stage is as (1).

$$I_d = \mu_n C_{ox} (W/L) [(V_{GS} - V_{DS})V_{DS} - V_{DS}^2/2] \quad (1)$$

Since the drain-to-source voltage of the control transistor (CT1) is zero, which makes  $V_{bs}=V_{GS}$ , the current  $I_d$  goes through ST1 can be written as (2)

$$I_d = \mu_n C_{ox} (W/L) [(V_{DS}^2/2 - V_{th}V_{DS})] \quad (2)$$

As the voltage level of virtual ground drops,  $V_{DS}$  over the ST1 drops and this makes the drain to source current of the sleep transistor (ST1)  $I_D$  drops quadratic manner. The dropping  $I_d$  decreases the voltage fluctuation on the ground and power net. Hence reduction in active power.

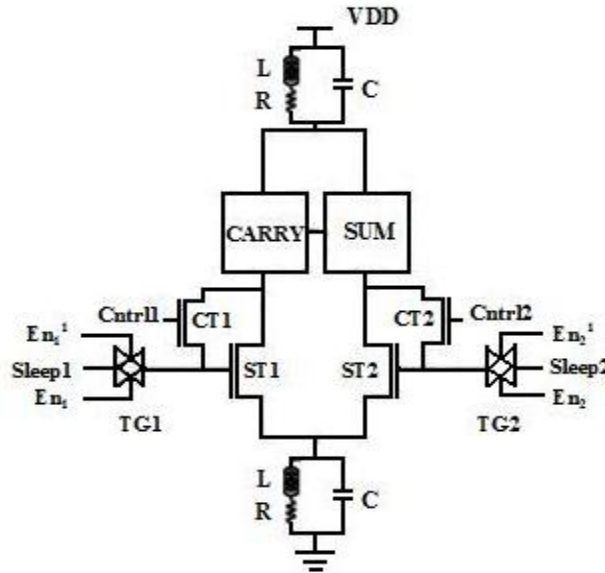


Fig. 6 Proposed novel technique

#### IV. GATE-DIFFUSION INPUT

The basic GDI cell is shown in Fig. 7 while the truth table is shown in Table I. It should be noted that the source of the PMOS in a GDI cell is not connected to VDD while the source of the NMOS in a GDI cell is not connected to GND. This feature gives the GDI cell two extra input pins to use which makes the GDI design more flexible than a usual CMOS design. However, this feature is also the major cause of its disadvantage: special CMOS process required. To be more specific, the GDI scheme requires twin-well CMOS or silicon on insulator (SOI) process to implement which is of course more expensive than the standard p-well CMOS process.

TABLE I. Truth Table of the Basic GDI Cell

N	P	G	Out	Function
'0'	B	A	$\overline{AB}$	F1
B	'1'	A	$\overline{A+B}$	F2
'1'	B	A	$A+B$	OR
B	'0'	A	$AB$	AND
C	B	A	$\overline{AB+AC}$	MUX
'0'	'1'	A	$\overline{A}$	NOT

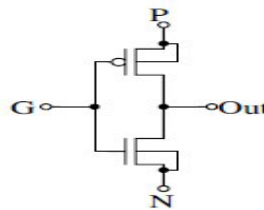


Fig. 7 Basic Gate-Diffusion Input Cell

**A. Conventional CMOS full adder**

The conventional CMOS logic gate full adder is shown as *Fig.8*, while the equation of a full adder is present as equation (3)-(6).

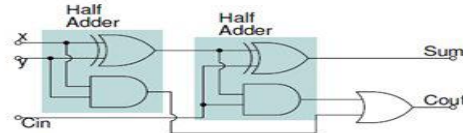
$$x+y+Cin = 2Cout+Sum \tag{3}$$

$$Cout = (y(x \oplus y)) + (Cin(x \oplus y)) \tag{4}$$

$$Sum = x \oplus y \oplus Cin \tag{5}$$

or

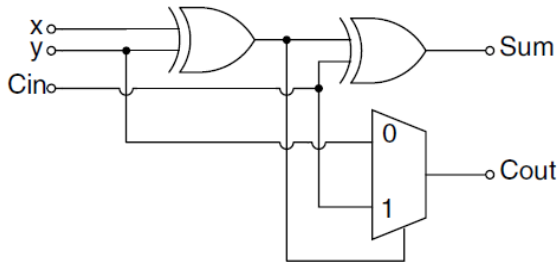
$$Sum = \overline{\overline{x \oplus y \oplus Cin}} \tag{6}$$



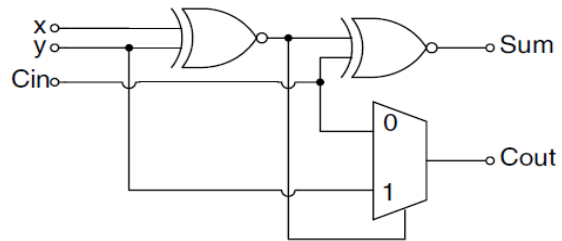
*Fig. 8* Conventional CMOS full adder

**B. XOR/XNOR based full adder**

According to equation (3) - (5), two XOR/XNOR based full-adders can be redesigned as shown in *Fig.9* (a) and (b), respectively. In addition, another two XOR/XNOR based full adders can be derived from equation (3), (4), and (6), respectively.



*Fig. 9* (a) XOR based full adder#1



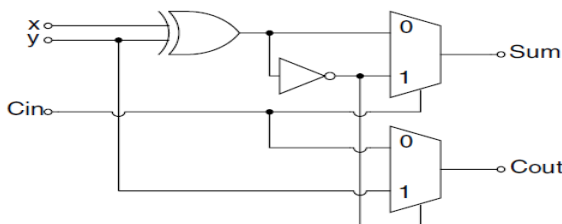
*Fig. 9* (b) XNOR based full adder#1

The schematics are shown as *Fig. 10*(a) and 10(b) [4]. Compare *Fig. 8* with *Fig. 9* and *Fig. 10*, it is obvious that the architecture of the adders in *Fig. 9* as well as *Fig. 10* are simpler than that of *Fig. 8*. However, the conventional CMOS XOR, XNOR and MUX gates require too many transistors to build and thus are not preferred choices. In this case, the XOR, XNOR, and MUX gates are redesigned using GDI scheme.

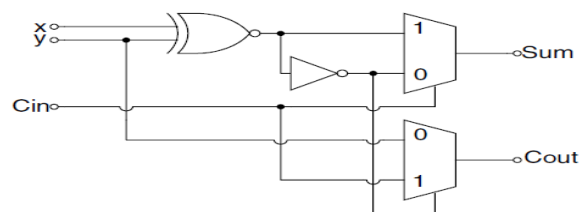
**V. REDESIGNED XOR/XNOR /MUX GATE**

**A. GDI XOR**

The GDI XOR gate is shown as *Fig. 11* where only 4 transistors are used. Compare the GDI XOR with its Conventional CMOS counterpart; it is obvious that GDI XOR gate requires fewer transistors.



*Fig.10* (a) XOR based full adder



*Fig.10* (b) XNOR based full adder

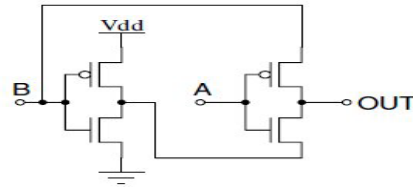


Fig.11.GDI XOR Gate

**B. GDI XNOR**

The GDI XNOR gate is shown as Fig. 12. It uses only 4 transistors as the GDI XOR gate does.

**C. GDI MUX**

Last but not least, a GDI MUX gate is implemented as shown in Fig. 13. The GDI MUX uses only 2 transistors.

**VI. THE PROPOSED GDI FULL ADDER**

According to the mentioned 4 different full adder architectures, we can redesign 4 different types of GDI based full adder architectures; we can redesign 4 different types of GDI based full adder.

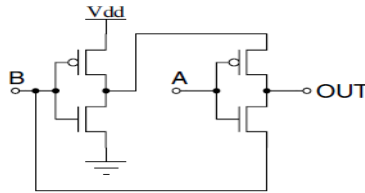


Fig. 12 GDI XNOR Gate

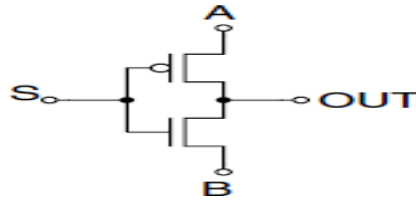


Fig. 13 GDI MUX Gate

These 4 adders are shown in Fig.14-17 respectively. It should be noted that all of the proposed full adders are 10-T based. Hence, the attempt to create 10-T based full adders is achieved.

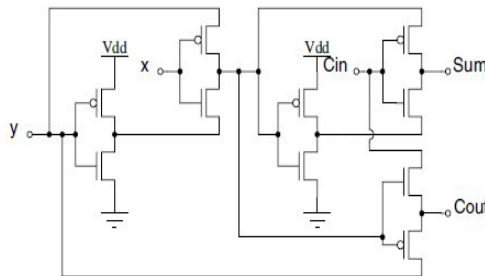


Fig.14 GDI XOR Full Adder#1(Based on Fig.8 (a))

**VII. COMPARISON**

In order to test the performance of the proposed GDI full adders, detailed comparisons are performed. In order to have a fair comparison, all the prior designs are recreated using the mentioned CMOS process. Besides, the transistor sizes of the prior designs are properly tuned to have optimal performance. Although we have performed variety comparisons under different temperature circumstances like 0°C, 25°C, 50°C, and 75°C.

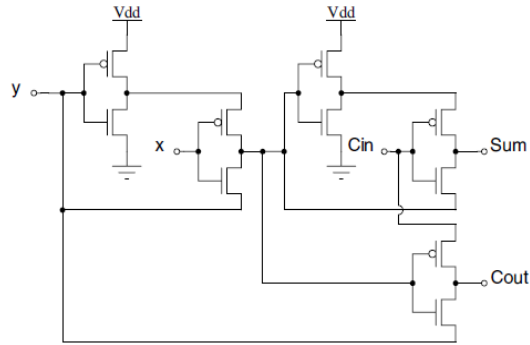


Fig. 15 GDI XNOR Full adder #1(Based on Fig.9 (b))

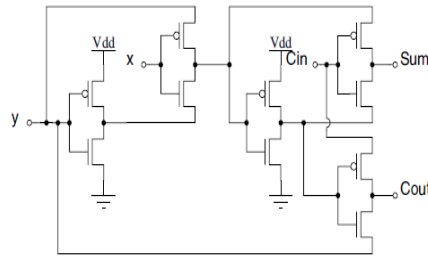


Fig. 16 GDI XOR Full adder #2(Based on Fig.10 (a))

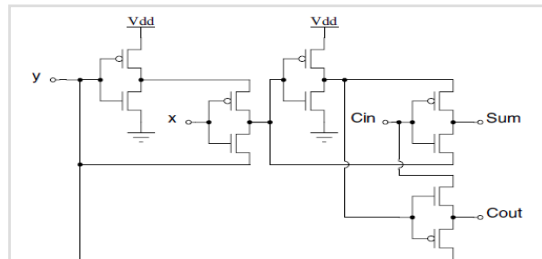
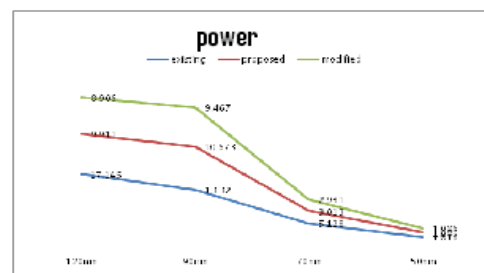
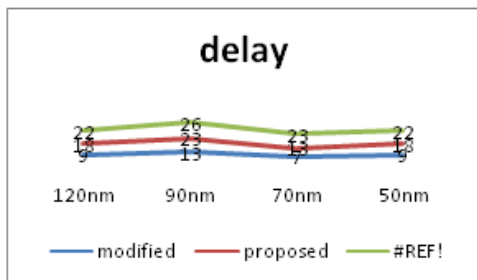


Fig. 17 GDI XNOR Full adder #2(Based on Fig. 10 (b))

It is very obvious that the proposed GDI XNOR FA #2 has the minimal Power Delay product. Besides, the proposed GDI XOR FA #1 also possesses very small Power Delay product. Hence, the proposed GDI XNOR FA #2 is the optimal 10-T based full adder design while the proposed GDI XOR FA #1 is also a preferable choice.

## VIII. RESULTS



## IX. CONCLUSION

Low leakage 1 bit full adder cells are proposed for mobile applications with low ground bounce noise. By using the proposed technique leakage power is reduced in (Design1) and (Design2) compare to the conventional adder cell (Base case). Ground bounce noise is reduced about 1.5 times and 3 times in Design1 and Design2 respectively compared to Base case. Further, using the proposed Novel technique the ground bounce noise is reduced to about 4.5 times in three designs (Base Case, Design1, and Design2) compared to without applying the technique. Area is reduced by 55.4% in (Design1), 72% (Design2) in comparison to the Base case. Active power reduction is reduced by 40.48% (Design1), 63.38% (Design2) in comparison to Base case. In this paper, 4 refined GDI based 10-T full adders are proposed. The proposed GDI design reduce active power by 80%.this design possesses the advantages of flexibility, less transistor counts, and can be realized using standard p-well CMOS process. In addition, a chip using the proposed designs is implemented. The chip is realized by Micro wind 90 nm, 0.12um and 70 nm CMOS technology. The comparison between our designs and prior works indicates that one of our designs does provide its advantages. In short, the proposed designs can be taken a better alternative.

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**S.Madhuri**, member IES obtained her B.Tech degree in Electronics and Control Engineering from JNTU, Hyderabad, is currently doing post graduation in Vaagdevi institute of Tech & Science, Proddatur, Kadapa (Dt), A.P with the specialization of VLSI, and also attended for 4 Faculty Development Programs and 1 Teaching Development Program.



**Avvaru Srinivasulu** received the B.Tech degree in Electronics and Control Engineering in 2006 and M.Tech degree in Instrumentation and Control Systems in 2008 from JNTU, Hyderabad. At present he is working as Assistant professor Dept.ECE. He published 2 papers in international journals and published one book on "FPGA implementation of Hopfield neural network for solving constraint satisfaction problems" by Lambert Academic Publishing, Germany. And also he attended for 8 Faculty Development Programs and 3 Teaching Development Programs.



**C.Mahammed Aslam** is the HOD, Dept.of Electronics and Control Engineering at Vaagdevi institute of Tech&Science, Proddatur. He obtained his B.Tech degree in ECE from Dr.Babasaheb Ambedkar Maratwada University, Aurangabad in 1997, M.Tech degree in Digital Electronics and Computer Science from JNTU, Hyderabad in 2007.He has 10 years of teaching experience and published 2 papers in international journals. He registered his Ph.D from JNTU-Anantapur in digital Image Processing.