

## Efficient dominating sets in ladder graphs

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**Abstract**—An independent set  $S$  of vertices in a graph is an efficient dominating set when each vertex not in  $S$  is adjacent to exactly one vertex in  $S$ . In this note, we prove that a ladder graph  $L_n$  has an efficient dominating set if and only if  $n$  be a multiple of 4. Also we determine the domination number of a ladder graph  $L_n$ .

**Keywords**—Efficient dominating set, Perfect code, Domination number, Ladder graph, Sphere packing

### I. INTRODUCTION

Let  $\Gamma = (V, E)$  be a finite undirected graph with no loops and multiple edges. We follow the terminology of [2] and [3]. Given  $S \subseteq V$ , let the open neighbourhood  $N(S)$  of  $S$  in  $\Gamma$  be the subset of vertices in  $V \setminus S$  adjacent to some vertex in  $S$ , and let the corresponding closed neighbourhood be  $N[S] = N(S) \cup S$ . A set  $S \subset V$  is a dominating set if  $N[S] = V$ , that is, every vertex in  $V \setminus S$  is adjacent to some vertex in  $S$ . The domination number  $\gamma(\Gamma)$  is the minimum cardinality of a dominating set in  $\Gamma$ . If the dominating set  $S$  is a stable set of  $\Gamma$ , then  $S$  is an independent dominating set. Also, when every vertex in  $V \setminus S$  is adjacent to exactly one vertex in  $S$ , then  $S$  is a perfect dominating set. A dominating set  $S$  which is both independent and perfect is an efficient dominating set. In what follows we may refer to an efficient dominating set as an  $E$ -set. Given the graph  $\Gamma = (V, E)$ , we call any subset  $C$  of  $V$  a code in  $\Gamma$ .

We say that  $C$  corrects  $t$  errors if and only if the sets  $S_c = \{u \mid u \in V, d(u, c) \leq t\}$  are pairwise disjoint. Moreover we call  $C$  a  $t$ -perfect code if and only if these sets form a partition of  $V$ . A  $t$ -perfect code  $C$  is called nontrivial if and only if  $t > 0$  and  $\text{card}(C) > 1$ .

$E$ -sets correspond to perfect 1-correcting codes in  $\Gamma$ , as treated by Biggs[1] and Kratochvil[4]. Equivalently, they provide a perfect packing of  $\Gamma$  by balls of radius 1. When  $\Gamma$  is  $r$ -regular, the so-called sphere packing condition  $|V| = (r + 1)|C|$  is trivially a necessary condition for  $C$  to be an  $E$ -set of  $\Gamma$ .

We need the following two results from [3]. We recall that the maximum degrees of vertices of  $\Gamma$  is denoted by  $\Delta(\Gamma)$ .

**Theorem 1.1** For any graph  $\Gamma$  of order  $n$ ,

$$\left\lceil \frac{n}{1 + \Delta(\Gamma)} \right\rceil \leq \gamma(\Gamma) \leq n - \Delta(\Gamma)$$

**Theorem 1.2** If  $\Gamma$  has an efficient dominating set, then the cardinality of any efficient dominating set equals the domination number  $\gamma(\Gamma)$ . In particular, all efficient dominating sets of  $\Gamma$  have the same cardinality.

### II. THE MAIN RESULTS

We use the following remark in the sequel

**Remark.** Suppose that  $n (> 1)$  be a natural number, the dihedral group  $D_{2n}$  of order  $2n$  is defined by the representation

$$D_{2n} = \langle x, y \mid x^n = y^2 = (xy)^2 = 1 \rangle.$$

Then the distinct elements of  $D_{2n}$  are as follows:

$$\{1, x, \dots, x^{n-1}, y, xy, \dots, x^{n-1}y\}.$$

The Cayley graph  $C(D_{2n}, X)$  of  $D_{2n}$  with respect to the generating set  $X = \{x, x^{-1}, y\}$  is a ladder graph  $L_n$ .

**Theorem 2.1** Let  $L_n$  be a ladder graph of order  $2n$ . Then the domination number of  $L_n$  is determined as follows:

$$\gamma(L_n) = \begin{cases} \frac{n}{2} & 4 \mid n \\ \frac{n}{2} + 1 & \text{if } 2 \mid n, 4 \nmid n \\ \left\lceil \frac{n}{2} \right\rceil + 1 & 2 \nmid n \end{cases}$$

**Proof.** If  $4|n$ , then  $n = 4k$  for some natural number  $k$ . Let  $S = \{1, x^4, x^8, \dots, x^{4k-4}, x^2y, x^6y, \dots, x^{4k-2}y\}$ . We show that  $S$  is a dominating set. It is easy to see that  $|S| = \frac{n}{2} = 2k$ . Since  $L_n \cong C(D_{2n}, X)$  then for all  $i$ ,  $2 \leq i \leq n-1$ , a vertex  $x^i$  is adjacent to three vertices  $x^{i+1}, x^{i-1}$  and  $x^i y$ . There are 3 cases.

If  $4|i$ , then  $x^i \in S$ .

If  $2|i$  and  $4 \nmid i$ , then  $x^i y \in S$ .

If  $2 \nmid i$ , then  $i = 4m-1$  or  $i = 4m+1$  for some natural number  $m$ .

If  $i = 4m-1$ , then  $x^{i+1} \in S$  and if  $i = 4m+1$ , then  $x^{i-1} \in S$ .

Similarly, for all  $i$ ,  $1 \leq i \leq n-1$ , a vertex  $x^i y$  is adjacent to some vertex in  $S$ . For three vertices  $1, x^{n-1}, y$ , we have  $1 \in S$  and  $y, x^{n-1}$  are adjacent to  $1$ . So  $S$  is a dominating set. Since  $L_n$  is a 3-regular graph, it follows from sphere packing condition that

$$\gamma(L_n) \geq \frac{n}{2}$$

Also  $S$  is a dominating set of  $L_n$  which  $|S| = \frac{n}{2}$ . Hence it follows that  $\gamma(L_n) = \frac{n}{2}$ ; ( $n = 4k$ ).

If  $n$  be odd, then we have 2 cases.

**Case1:**  $n = 4k+1$  for some natural number  $k$ .

In this case, let  $S = \{1, x^4, x^8, \dots, x^{4k}, x^2y, x^6y, \dots, x^{4k-2}y\}$ , then

$$|S| = 2k + 1 = \left\lfloor \frac{4k+1}{2} \right\rfloor + 1.$$

It is easy to see that  $S$  is a dominating set, and by theorem 1.1,  $\left\lfloor \frac{n}{2} \right\rfloor + 1 \leq \gamma(L_n) \leq |S| = \left\lfloor \frac{n}{2} \right\rfloor + 1$ .

Therefore,  $\gamma(L_n) = \left\lfloor \frac{n}{2} \right\rfloor + 1$ ; ( $n=4k+1$ ).

**Case2:**  $n = 4k-1$  for some natural number  $k$ .

Let  $S = \{1, x^4, x^8, \dots, x^{4k-4}, x^2y, x^6y, \dots, x^{4k-2}y\}$ , then

$$|S| = 2k = \left\lfloor \frac{4k-1}{2} \right\rfloor + 1 = \left\lfloor \frac{n}{2} \right\rfloor + 1.$$

It is similar to case1. So  $\gamma(L_n) = \left\lfloor \frac{n}{2} \right\rfloor + 1$ ; ( $n = 4k-1$ ).

If  $2|n$  and  $4 \nmid n$ , then  $n = 2k$  such that  $k$  is an odd number. Let

$$S = \{1, x^4, x^8, \dots, x^{2k-2}, x^2y, x^6y, \dots, x^{2k-4}y, x^{2k-1}\}.$$

It is easy to see that  $S$  is a dominating set and  $|S| = k + 1 = \frac{n}{2} + 1$ .

By theorem 1.1,

$$\frac{n}{2} \leq \gamma(L_n) \leq |S| = \frac{n}{2} + 1.$$

We show that the domination number  $\gamma(L_n) \neq \frac{n}{2}$ . It follows from the proof of theorem 1.1, that  $\gamma(\Gamma) = \frac{n}{1+\Delta(\Gamma)}$  if and only if  $\Gamma$  has a  $\gamma$ -set (dominating set of minimum cardinality)  $S$  such that  $N[u] \cap N[v] = \emptyset$  for all  $u, v \in S$  and  $|N[v]| = \Delta(\Gamma)$  for all  $v \in S$ . If  $\gamma(L_n) = \frac{n}{2}$ , then  $L_n$  has a  $\gamma$ -set  $S$  such that  $N[u] \cap N[v] = \emptyset$  for all  $u, v \in S$ . Since  $N[S] = V(L_n)$ , therefore

$$3k = |N[S]| = |V(L_n)| = 2n = 4k$$

is a contradiction. So  $\gamma(L_n) \neq \frac{n}{2}$ . ■

**Remark.** If the ladder graph  $L_n$ , which is a 3-regular graph has an efficient dominating set  $S$ , then by the sphere packing condition we have  $|S| = \frac{n}{2}$ . Therefore, only the ladder graph with even order has an efficient dominating set.

**Theorem 2.2** The ladder graph  $L_n$  of order  $2n$ , has an efficient dominating set if and only if  $4|n$ .

**Proof.** By remark and this fact that  $|S| = \gamma(L_n)$ , the ladder graph has an E-set if  $4|n$ . Also if  $4|n$ , then we show that the set

$$S = \{1, x^4, x^8, \dots, x^{4k-4}, x^2y, x^6y, \dots, x^{4k-2}y\}$$

is a dominating set. It is easy to see that  $S$  is an independent and perfect set. Therefore  $S$  is an E-set of ladder graph  $L_n$ . ■

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