Analysis of Improved DV-Distance Algorithm for Distributed Localization in WSNs.

1Sudha H. Thimmaiah and 2G. Mahadevan
1Research Scholar, Dept. of ECE, PRIST University, Thanjavur, Tamil Nadu, India.
2Principal, Dept. of CSE, Annai College of Engineering, Kumbakonam, Tamil Nadu, India.
Corresponding Author: sudhathimmaiah@yahoo.com

ABSTRACT: Node localization is one of the important supporting technologies in Wireless Sensor Networks. The process of node localization depends on the accuracy of determining the position of the nodes in the required environment. Most of the applications of WSNs require the knowledge of the location of the sensors. The main task of a wireless sensor node is to sense and collect data from a certain domain, process the data and transmit the same to the anchor node. In this paper, both the DV-distance localization algorithm and the improved version of the same is analysed for localization errors. In the DV-distance algorithm, the cumulative distance is obtained by the use of distance vector routing whereas, in the improved DV-distance algorithm hops are introduced between anchor and unknown nodes to determine the distance. Simulation results show that, localization accuracy of the improved algorithm is better than the DV-distance algorithm.


Date of Submission: 10-02-2018
Date of Acceptance: 26-02-2018

I. Introduction

Wireless Sensor Networks is one of the recent networking fields which have gained significance in several key research areas. A wireless sensor network is composed of a large number of sensor nodes densely deployed in the sensor network domain. The sensor nodes in the network domain must be able to sense and collect data, process the data and transmit the data to the anchor or beacon node in most applications. Wireless Sensor Networks have wide application prospects including search and rescue, disaster relief, target tracking and smart environments. Hence, location awareness is a key factor for the success of the applications. According to the different information needed in the process of node localization, positioning algorithm can be classified into two kinds: Range-based and Range-free [1]. Dragos Niculescu in Rutgers University, who depended on the distance vector routing and GPS principles, proposed a series of distributed positioning algorithms which is well known as Ad Hoc Positioning System – APS [2]. It includes six algorithms: DV-Hop, DV-Distance, Euclidean, DV-Coordinate, DV-Bearing and DV-Radial. Amongst these, the most popular and successful algorithm is the DV-Hop and DV-Distance localization algorithms.

II. Review Of DV-Distance Algorithm

The DV-Distance is similar to the DV-Hop, but with the difference that distance between neighbouring nodes is measured using Received Signal Strength Indication (RSSI) and the data is propagated in meters rather than hops. The cumulative distance is then calculated by using the method similar to the distance vector routing. DV-Hop localization approach is composed of three phases:

- In the first phase, a typical distance vector routing mechanism is employed. Beacons flood their location information throughout the network with the initial hop-count of zero. Each node that relays the message increases the hop-count by one. After the flooding procedure, every node can obtain the minimum hop count to each beacon.

- In the second phase, each beacon, after obtaining the position and hop-count information to all other beacons, estimates the average distance per hop. For example, beacon \( i \) calculated the average distance per hop, called the hop size HS, using the formula,

\[
HS_i = \sum_j (x_i - x_j)^2 + (y_i - y_j)^2 / \sum h_j
\]

Where \((x_i, y_i)\) and \((x_j, y_j)\) are the coordinates of beacons \(i\) and \(j\) respectively, and \(h_j\) is the hop-count value from beacon \(i\) to beacon \(j\).
In the last phase, before conducting the self-localization, each sensor estimates the distance to each beacon based on its hop-count and the hop-size to this beacon. For example, sensor \(k\), can get the distance \(dkj\), the distance from sensor \(k\) to beacon \(j\) by using, \(dkj = hj \times HSj\). After obtaining all the distance information, each sensor conducts the triangulation or maximum likelihood estimation to estimate its own location. This localization algorithm will only work well in isotropic networks [3].

The DV-Distance localization algorithm in [6] uses the method of distance vector routing to obtain the cumulative distance. When the unknown nodes get three or more cumulative distances from anchors, then the trilateration method is used to calculate the localization. The DV-distance localization algorithm can be divided into three steps [7].

- In the first step, each anchor node broadcasts a beacon to be flooded through the network containing the anchors localization with node ID and RSSI. Each receiving node calculates the distance between its adjacent nodes based on the RSSI theoretical model. Then they will count the cumulative distance or the polyline-distance \((dis_{i,j})\) from themselves to anchors that they can receive by cumulating the polyline-distance between hop to hop.
- In the second step, once an anchor node \(i\) gets the polyline-distance to another anchor node \(j\), it calculates an correction ratio \((cor_{i,j})\) of polyline-distance \((dis_{i,j})\) to Straight-line \((Dis_{i,j})\) between the two anchors \(i,j\) based on (1), which is then flooded to the entire network as the network correction. When an unknown node \(k\) accepts the information from one anchor node \(j\), it uses the polyline-distance \((dis_{i,j})\) from anchor nodes to compute the amended distance \((DisCorr_{k,j})\) between itself and the anchor node as the following Equation (2)

\[
cor_{i,j} = \frac{dis_{i,j}}{Dis_{i,j}} \quad (1)
\]

\[
DisCorr_{k,j} = \frac{dis_{k,j}}{correction} \quad (2)
\]

- In the last step, when the unknown nodes get three or more distances from anchors, and then unknown nodes calculate their localization by trilateration or maximum likelihood estimation.

DV-Distance replaces the hops in DV-Hop with the polyline-distance which is calculated based on the RSSI model. Thus, it reduces the error due to consider that the rough distance of each hop is equal. The implementation of the algorithm is easy, but, with the ranging error increasing, the localization error also sharply increases [4]. The DV-Distance localization algorithm for localization is used when all the nodes in the network have the ranging capability.

### III. Review Of Improved DV-Distance Algorithm

The improved DV-Distance algorithm takes into account most of the jumping segment of the unknown node and anchors from its nearest anchor node to other anchors is overlap. So, when the unknown node selects the distance correction parameter \((Correction)\), the unknown node selects the original distance correction parameter which has been broadcasted to the networks, if the anchor is the closest node to this unknown node. Otherwise, the unknown nodes use corrections of its nearest anchor node to the other anchor nodes, instead of the unknown node to the other anchor nodes. The proposed algorithm, only requests the anchors have the capability to save the Corrections [5].

- In the first step, the proposed algorithm is similar with DV-Distance. The difference is that it uses the method similar with the distance vector routing to get the cumulative distance and the cumulative hops.
- In the second step, once an anchor gets the polyline distance to other anchors, it calculates the correction with the following Equation (3)

\[
Correction_{i,j} = \frac{(dis_{i,j} - Dis_{i,j})/Hop_{i,j}}{Dis_{i,j}} \quad (3)
\]

Where, \(dis_{i,j}\) is the polyline-distance of anchor \(i\) and \(j\), \(Dis_{i,j}\) is the beeline-distance of anchor \(i\) and \(j\), \(Hop_{i,j}\) is the hops of anchor \(i\) and \(j\). \(Correction_{i,j}\) is the distance correction parameter.
Then, it saves the Correction and the anchor ID to its established list in the memory, and then broadcasts the Correction which is the correction parameter form the nearest anchor node to the network, using controlled flooding. This scheme could ensure that the most nodes receive the Correction from beacon node that has the least distance. When the unknown node gets the Correction from an anchor node, and then it saves this Correction as the correction parameter between this unknown node and its nearest anchor node. Then, the unknown node returns its node ID and other anchor ID which has been saved in the Memory list of this unknown node, and then the answer anchor node sends the corrections which the unknown node need back to the unknown node. The unknown node gets the corrections to all the anchors that its ID has been saved. Then, it calculates the corrected distance with the following equation (4)

\[
DisCorr_{k,i} = \frac{dis_{k,i}}{1 + \text{Correction}_{k,i} \times \text{Hop}_{k,i}}
\]

Where, \(dis_{k,i}\) is the polyline-distance of anchor i and unknown node k, \(\text{Correction}_{k,i}\) is the correction of anchor i and unknown node j, \(\text{DisCorr}_{k,i}\) is the cumulative-distance of anchor i and unknown node.

In the last step, implementing the trilateration positioning gets the coordinates of all unknown nodes.

IV. Methodology / Implementation

Wireless Sensor Networks consists of three kinds of nodes – target node, unknown node and anchor node. The number of hops and polyline distance of each unknown node should be found out with respect to anchor nodes, the number of hops and polyline distance between target node and anchor node should be determined. The correction is done for all the parameters using improved DV-Distance method and trilateration to get the localization co-ordinates [6, 7]. MATLAB 9.0 is used to simulate the improved DV-Distance algorithm. The proposed algorithm is divided into three modules: (1) WSN Deployment (2) Improved DV-Distance method (3) Trilateration

(1) WSN Deployment

In this module, the specification of anchor nodes and unknown nodes is given as the input, the co-ordinates of each node are generated randomly and the communication range between one node and the other is defined to get the set of nodes within the communication range to form the connected WSN and represented in figures 1 and 2.

Input: The number of anchor and unknown nodes
Output: Connected WSN

Random Distribution of Sensor Nodes

Fig 1. 4 anchor sensors and 100 unknown sensors randomly deployed in 80x80 unit square area.

Sensor Nodes Communicating Within a Sensing Range
Analysis of Improved DV-Distance Algorithm for Distributed Localization in WSNs.

Fig 2. The Sensor nodes communicating within the sensing range.

(2) Improved DV-Distance method:
Improved DV-Distance algorithm is applied to the connected WSN and represented in figure 3.
Input: Connected network, polyline distance, number of hops, beeline distance.
Output: the corrected distance of each unknown nodes to the nearest anchor nodes.

(3) Trilateration:
The trilateration method is applied to obtain the estimated position of unknown nodes as shown in figure 3.
Input: the corrected distance of each unknown nodes to the nearest anchor nodes.
Output: The estimated localization co-ordinates of unknown nodes.

Localization of all unknown nodes

Fig 3. Plot between each node and the estimated position of x-coordinate and y-coordinate values after the improved DV-Distance and trilateration method is applied for the estimation position of unknown nodes.

Performance of localization
Figure 4 represents a plot of node numbers on the X – axis and the mean absolute error in the Y – axis and hence the performance of localization.

Fig 4. Errors in the corresponding estimated position of unknown nodes.
V. Results And Discussion

The simulation results showed the deployment of random nodes, their connections within the communicating range, the estimation of localization with improved DV-Distance algorithm and trilateration method [8,9]. The obtained localization of unknown nodes shows very little error compared to the exact localization values. The mean absolute error in the localization of unknown nodes is almost an average throughout with only a few peaks at some nodes [10, 11].

VI. Conclusion

Hence, the estimation of the location accuracy is improved and thus verified. Further improvement in the location accuracy can be increased with more computing complexity and storage capacity of anchor nodes.

References